

Formal Contracts for Runtime Verification Support in the Ada Programming Language



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Motivation

- Static Verification is not sufficient to cope with many of the challenges of modern and future generation real-time embedded systems:
 - state-explosion problem of model-checking;
 - limited automation in deductive reasoning, even with recent advances in SAT and SMT solvers.
- Most of the data important to certify a real-time embedded system is related to extra-functional properties:
 - Duration of tasks;
 - Energy consumption;
 - Temperature management;
 - Other cyber-physical properties.
- Unfortunately, most of the extra functional data is only available and verifiable during execution time.

Runtime Verification

- Runtime Verification is the discipline that studies formal theories and that proposes methods to generate monitors capable of observing and verifying formal specification during execution time:
 1. Formal specifications determine the property of interest that must be verified;
 2. Monitors are generated from that specification and are instrumented into the system.
- Typical contracts establish properties about the program that are verified via static approaches
- Runtime Verification behavior should follow the same principles:
 - Users define contracts about properties that he wishes to see verified upon execution;
 - The system is responsible for generating the monitors from those contracts.

Ada 2012 and Contracts

- Contracts enhance trust in the system by establishing a compromise between requirements and implementation
- Ada 2012 provides a sub-language for specifying contracts:
 - Checked at runtime via asserts, or;
 - Statically verified using the SPARK toolset.
- Contract language provides the ideal environment to specify properties that we need to be checked upon run-time (e.g., timed behavior of tasks)
- Runtime Verification contracts can be pre-processed to generate the monitors, and afterward removed, thus preserving the standard Ada 2012 contracts

Underlying Architecture

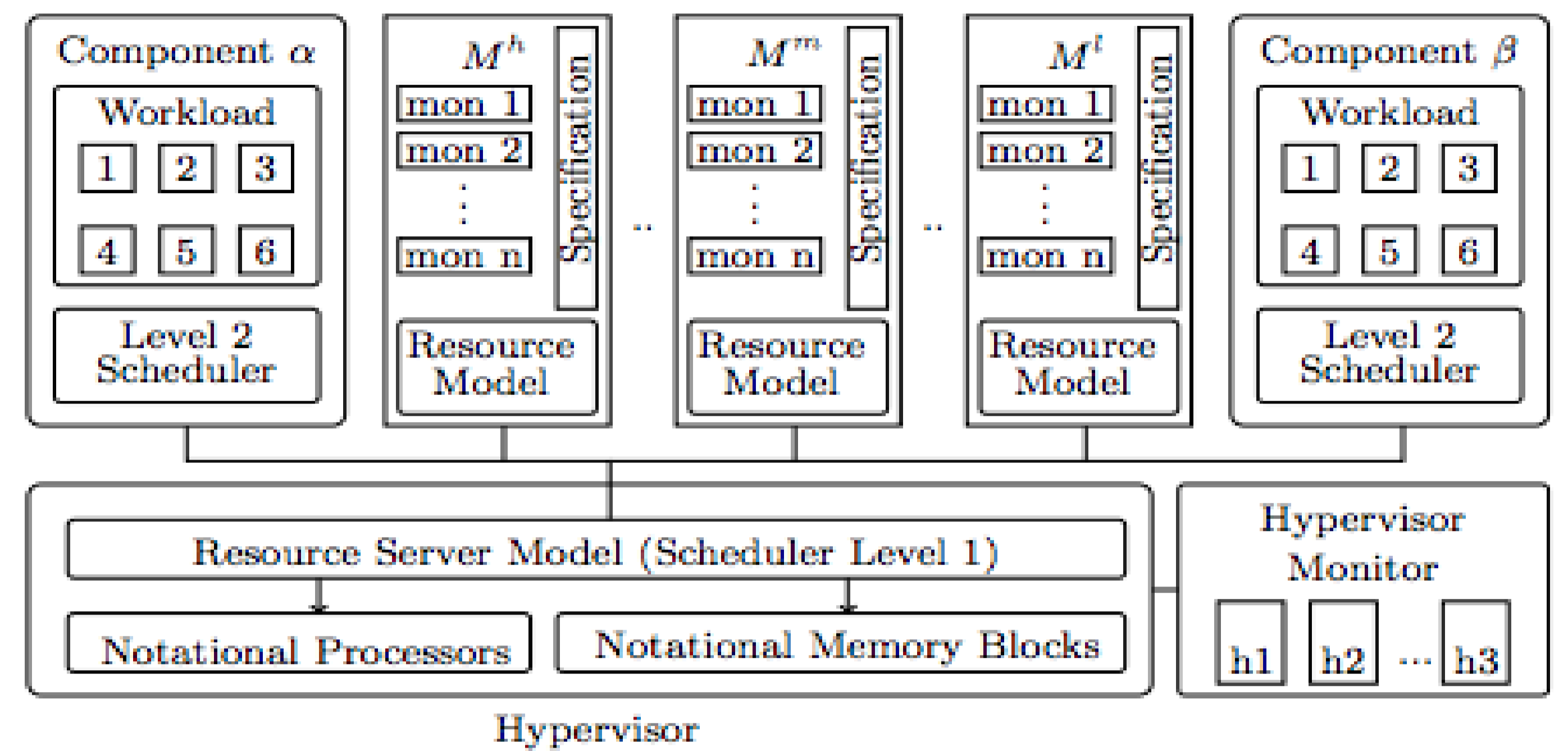


Fig. 1: Component-based Monitoring Architecture (CMA)

Pluggable Formal Theories

Timed Regular Expressions

```
task type T.Simulation (period: integer; deadline: integer)
with
  Monitor_Mode => Event_Triggered,
  Monitor_Case => ( RMTLD ,
    T.Simulation'Event(Task_Release) next implies
    duration[T.Simulation'Time(period)]
    T.Simulation'Event(ANY) < T.Simulation'Time(wcet)
  );
```

Metric Temporal Logic with Durations

```
protected type Protected_Environment
with
  Monitor_Mode => Time_Triggered
  Monitor_Case => ( TRE,
    ( Protected_Environment.read_CH4'Event(pre) .
    <(Protected_Environment'Event(ANY))>[0..20] .
    Protected_Environment.read_CO2'Event(post)) ) ,
is
  function read_CO2 return CO2_Level_State;
  function read_CH4 return CH4_Level_State;
  function read_Air_Flow return Air_Exhaust_State;
  function read_WaterPipe_Flow return WaterPipe_Flow_State;
end;
```

References

- [1] Pedro, A., Pereira, D., Pinho, L.M., Pinto, J.S. "Towards a Runtime Verification Framework for the Ada Programming Language". Reliable Software Technologies - Ada Europe 2014, LNCS 8454, pp. 58-73, Paris, France, 2014.
- [2] Pedro, A., Pereira, D., Pinho, L.M., Pinto, J.S. "A Compositional Monitoring Framework for Hard Real-Time Systems". NASA Formal Methods Symposium 2014, LNCS 8430, pp. 16-30, Houston, Texas, USA, 2014.